

BIG DATA FOR SMALL DOLLARS.

NEIL STEVENSON 11:55, 25TH JUNE

ABOUT ME - NEIL STEVENSON

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- Solution architect for Hazelcast
- Started in IT in 1989
- Has maintained programs written before he was born
- Fond of coffee , beer, and coffee
- Mainly a Java person, some GoLang
- Remembers the launch of C++
- Knows what IEFBR14 is



BIG DATA

Who remembers the "Y2K Problem"?

- Data records looked like "SW1V1EQ 1155180625".
- POSTCODE, byte[8]
- TIME, byte[4]
- DAY, byte[6]
 - This was BIG data! We could not afford 8 bytes for day



BIG DATA

- BIG DATA == Data we cannot afford to store
- Storage costs money
 - **\$\$\$\$\$**
 - £££££
- Storage is cheaper and bigger than Y2K days
 - But data is bigger too, increasing at a faster rate, so the problem isn't going away



BIG DATA

- BIG DATA == Data we cannot afford to store
- Storage costs time
 - Store then compute, results arrive too late, for some applications
 - Even with in-memory storage!
 - So we need in-memory computing!



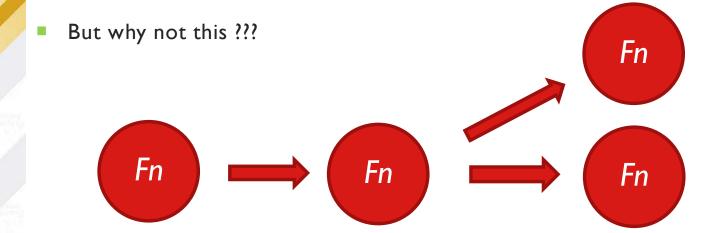
UNIX

- This is a Unix command "ls | grep neil | wc -1".
 - "1s" == no input, output is list of files
 - Discrete, output is produced then command ends
 - "grep neil" == filter for input containing the word *neil*, output the matches
 - Continuous, output produced as input arrives
 - "wc -1" == count the input, output the count
 - Discrete, output produced when input exhausted
- It's a simple chain of processing, no intermediate storage



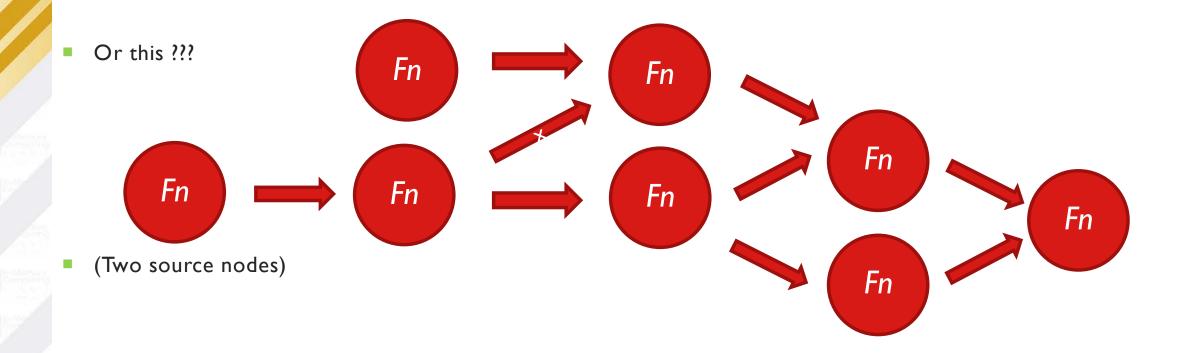
Really it's this:



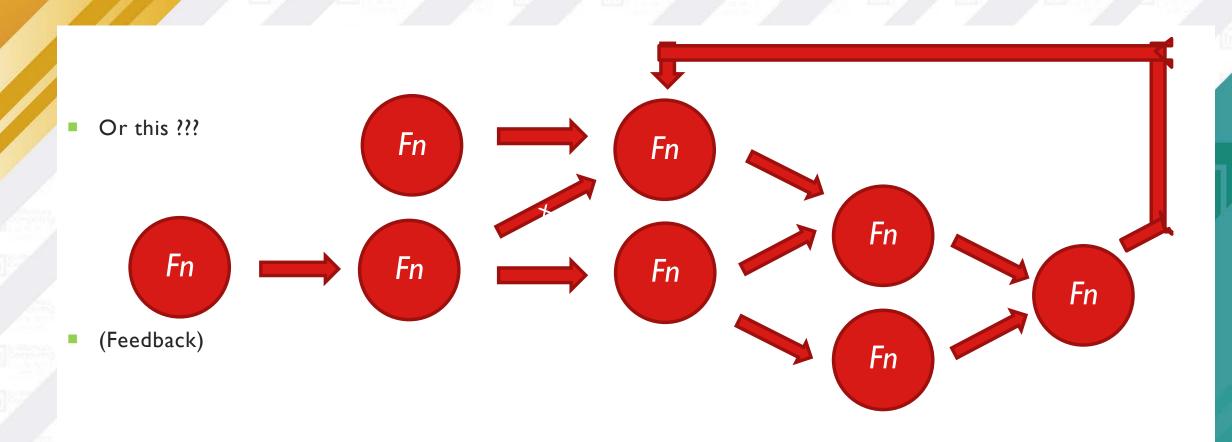


The "tee" command ??











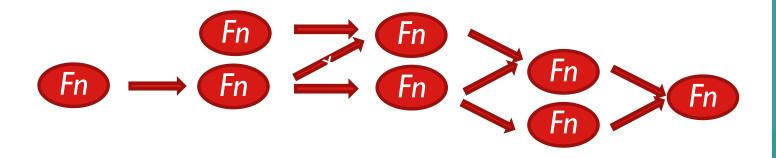


- Java based
- Open source
- Apache 2 licensed
- Distributed Streaming Analytics Engine
- Integrates trivially with Hazelcast IMDG
- Really good, says Neil that works for Hazelcast ©





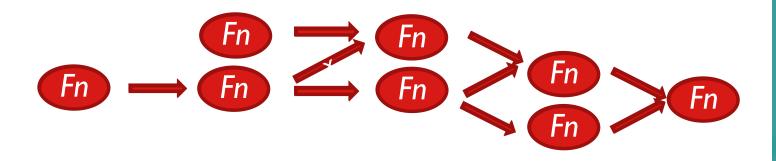
- Based around <u>acvelic graphs</u>.
 - No feedback loops







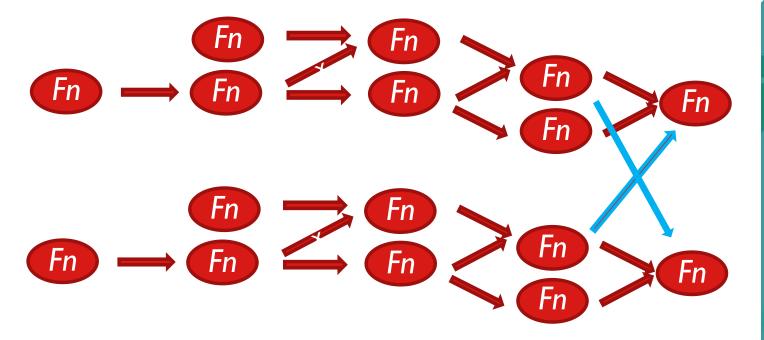
- But <u>distributed</u> acyclic graphs.
 - If you have 2 CPUs, run it twice
 - Different JVM or same JVM







- But **distributed** acyclic graphs.
 - If you have 2 CPUs, run it twice
 - Different JVM or same JVM
 - Data can cross instances





THE UBIQUITOUS "WORD COUNT"

```
Pipeline pipeline = Pipeline.create();

pipeline.drawFrom(Sources.<Integer, String>map("hamlet"))

flatMap(entry -> Traversers.traverseArray(Pattern.compile("\\W+").split(entry.getValue())))
.map(String::toLowerCase)
.filter(s -> s.length() > 3)
.groupingKey(DistributedFunctions.wholeItem())
.aggregate(AggregateOperations.counting())
drainTo(Sinks.map("count"));
```

Quiz time: Can you spot the mistake ?????



THE UBIQUITOUS "WORD COUNT"

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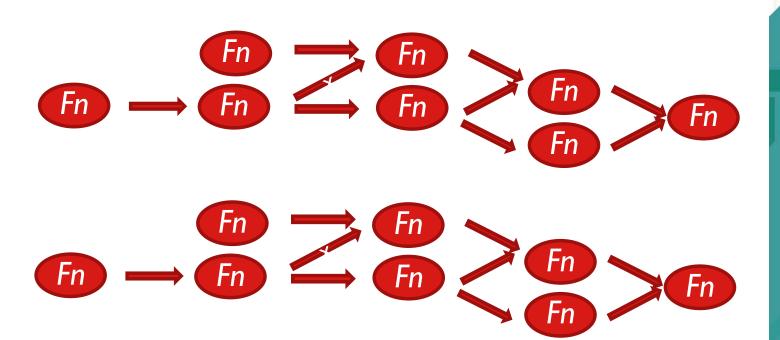
Answer: Filter on length is more efficient if it precedes "toLowerCase()". Performance cost!!! Not trivial



Data ingest is in parallel

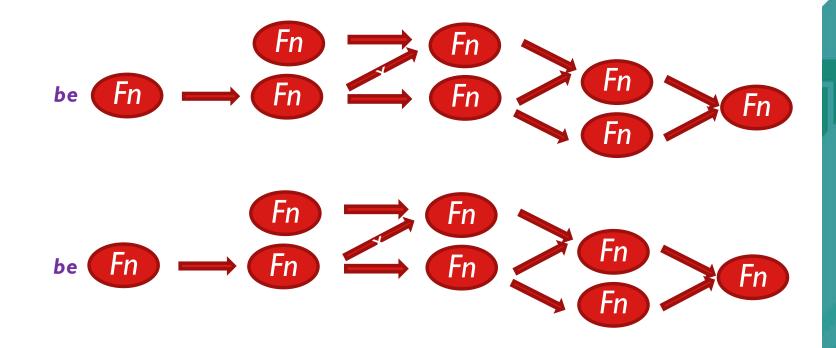
To be

Or not to be



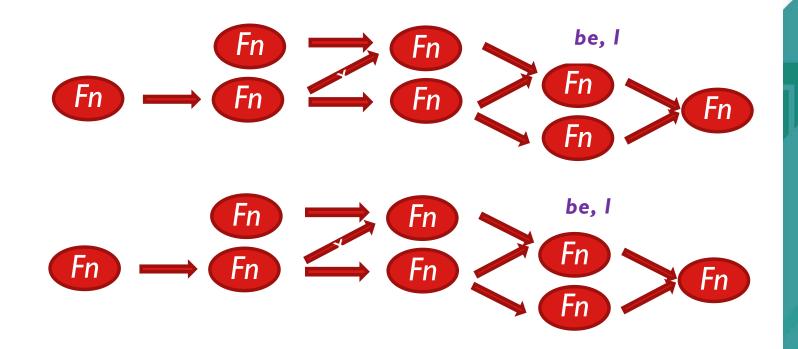


Data ingest is in parallel



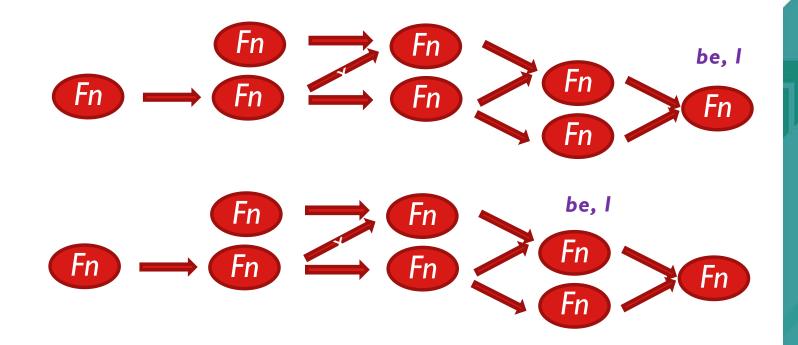


- Data ingest is in parallel
- Data egest is in parallel
- ..if you want



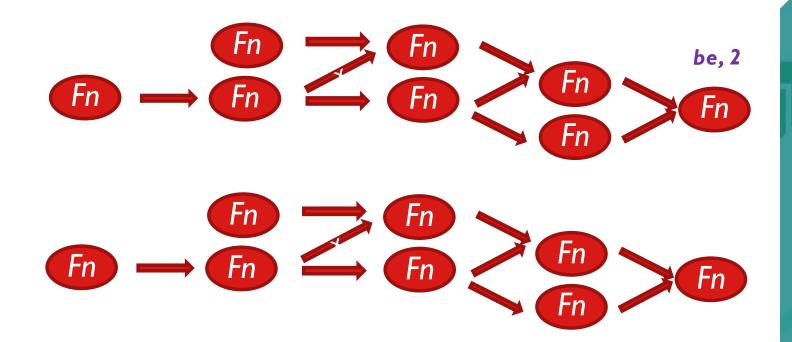


- Data ingest is in parallel
- Data egest is in parallel
- ..if you want





- Data ingest is in parallel
- Data egest is in parallel
- ..if you want





MEANWHILE

Ok, we have fast streaming processing....

Next we need some data, BIG data



WHAT IS BIG

- Superbowl 2018
 - Eagles v Patriots, 103.4 million viewers
 - https://www.cbsnews.com/news/super-bowl-lii-tv-ratings/
- Superbowl 2018 Half-Time Show
 - Justin Timberlake, 106.6 million viewers
 - http://money.cnn.com/2018/02/05/media/super-bowl-ratings/index.html
- World Cup 2014
 - Argentina v Germany final, 1.013 billion viewers
 - https://www.fifa.com/worldcup/news/2014-fifa-world-cuptm-reached-3-2-billion-viewers-one-billion-watched--2745519



THE 2014 WORLD CUP FINAL

- The final had 280 MILLION ONLINE viewers
- Many of these have Twitter accounts and will be tweeting
 - 674 million tweets about the final, before, during and after
 - Peak at 618,000 a minute (when Germany scored)



SO....

- Twitter is already storing the tweets, but we'd like to analyse them
- We want to do sentiment analysis
 - Who do the fans think will win before the game starts?
 - Who do the fans think will win while the game is in progress?
 - Why do we want to do this?
 - Place a bet on the winner! Make SMALL DOLLARS



THE PIPELINE

- Twitter firehose, tweets by hashtag
 <= could be parallel input across multiple JVMs</p>
- Filter out if not ASCII
- Enrich by locating a named team
- Filter out if no team named
- Filter out if team named not playing in this game
- Enrich with sentiment
- Increment running totals <= possible contention point, unless routing is used</p>



THE PIPELINE

- Twitter firehose, tweets by hashtag
- Filter out if not ASCII
- Enrich by locating a named team
- | Filter out if no team named <= Route here on team name
- Filter out if team named not playing in this game
- Enrich with sentiment <= Or is here better ?</p>
- Increment running totals



DEMO TIME

- Let's see code
 - java -jar target/worldcup-0.0.1-SNAPSHOT.jar

Uruguay v Russia is today at 3pm



DEMO TIME

Join in!!!

- Uruguay v Russia is today at 3pm
 - Hashtag "#URURUS"



DOES THIS WORK?

- No
- Or not yet, the business logic is too naïve
- But the idea is sound
- Download the code and fix it yourself ©



DOES THIS WORK?

- Some successes!
- Argentina v Croatia, after 18 minutes the sentiment at 0-0 was Argentina to lose. Final score 0-3
- Iran v Spain, at half-time and 0-0 the sentiment was for draw. Final score was 0-1, but Iran had a goal disallowed
- Uruguay v Saudi Arabia, at half-time and 0-0 the sentiment was for Uruguay. Final score was 1-0.
- But most of the others were wrong, so I'm not betting any money on the "predictions"



SUMMARY

- Stream processing == processing before storage
 - Someone else has stored already, eg. an IMDG
 - Can't afford cost of storage
 - Can't afford time for storage
- Distributed pipeline is a way to think about processing as a chain of simpler steps
 - Can benefit from machine parallisation



SUMMARY

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- https://github.com/neilstevenson/worldcup
 - You will need your own Twitter credentials

• Questions ?

